Common MikroTik OSPF mistakes and how to avoid them

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About me

Lorenzo Busatti

- Founder of Grifonline S.r.l. [ISP] 1997
- A user of MikroTik since 2006
- Founder of Linkwave [WISP] 2006
- MikroTik Trainer 2010
- Member of RIPE, AMS-IX, MIX-IT









About me









About me

Founder (2016) of the



High Quality Training Classes



About me

One of the founders (2017) of the Riga Bootcamp!





The Riga Bootcamp





The Riga Bootcamp





The Riga Bootcamp









The Schedule



THE FULL SCHEDULE





Dedicated to Max



Abstract

Installing and configuring OSPF is about understanding more than just how to configure RouterOS.

Avoiding common mistakes and assumptions made when deploying OSPF with RouterOS, leads to a higher quality of installation.

We will be covering some of the pitfalls we've seen and show how to avoid them.



Abstract

Delivering MTCRE and MTCINE Training Classes makes me aware of the lack of knowledge of the protocol and many misunderstandings by most of the technicians.



Abstract

The points we will focus on today:

- OSPF basics
- Common OSPF mistakes
- OSPF Extra tips



The OSPF protocol Open Shortest Path First

OSPF calculates the **shortest route** to a destination through the network based on the **Dijkstra's algorithm**.

If the network topology **changes** then the routing tables will also be **recalculated again**.



The OSPF protocol Open Shortest Path First

- Is a routing protocol for IP networks.
- It uses a link state routing (LSR) algorithm.
- It falls into the group of IGPs (interior gateway protocols).
- Is within a single AS (autonomous system).
- It is defined as OSPF Version 2 for IPv4.
- The updates for IPv6 are specified as OSPF Version 3.



The Dijkstra's algorithm

 Also known as the "shortest path first (SPF)"

 Is an algorithm for finding the shortest paths between nodes in a graph (an abstract data type)



The Dijkstra's Algorithm

Shortest path

 \checkmark means \checkmark

the "cheapest" path

 \checkmark means \checkmark

The **sum** of the cost of the **output** interface of **each router** for **the full path**



It's quite easy to implement the OSPF protocol in a network:

- 1) Put ip addresses in the same subnet between the routers;
- 2) Paste this code in each one:

/routing ospf network add
network=0.0.0.0/0 area=backbone





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Your network will start to run the OSPF protocol and now you start to think that you are an OSPF engineer!



- Network Type?
- Priority?
- DR?
- BDR?

I don't need to know these f***g things, by default everything is working!



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But when this network starts to grow you will not be smiling anymore.

That's why you need a better knowledge of how OSPF works!

I will not cover all the aspects of the OSPF and, of course, I will not deliver a MTCRE training class in 20 minutes ⁽²⁾

We will see some of the most interesting things, tips and mistakes that we can do using the OSPF protocol.



The OSPF Router's POV (Point Of View)



In OSPF networks there is <u>**no</u></u> "main" or "central" router that knows the "topology" of the network and the shortest path between them.</u>**

An **OSPF Router thinks they are the "core"** of their network, regardless of their real 'position' in your network.







Then each router, individually, calculates all the paths and their costs, for each known destination.

Then chooses the "shortest" (or cheapest) one.

From <u>his</u> own point of view.

(we will not show ALL the possible combinations)









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As I showed the paths **from R1 to R3**, that are only **in one direction**, let's now look at the paths **back from R3 to R1**.









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CADEM





One of the common mistakes is to check the paths from one side only.

That's why **you should look** at the paths **from both sides**.


The role of the backbone area



The backbone area (area-id=0.0.0.0) forms the core of an OSPF network.





The backbone is responsible for **distributing routing information** between other **non-backbone** areas.



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Each non-backbone area must be directly **connected to the backbone area** (**directly** or through a 'virtual link')



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Each non-backbone area must be directly **connected to the backbone area** (directly or through a '**virtual link**')



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By the previous definitions you can create no more than two "layers" of areas on top the backbone one.



The structure viewed from the side

Virtual Links





Or like a flower ⁽²⁾ (viewed from top)



Common OSPF mistakes



Common OSPF mistakes



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Some common thoughts about OSPF:

• "OSPF is weird."

- "It works when it wants to."
- "When it doesn't want to work it doesn't work."

Let's take a look

OSPF Network Types

From the wiki.mikrotik.com:

network-type (*broadcast I nbma I point-to-point I ptmp*; Default: **broadcast**)

Broadcast and the *nbma* types elect a **designated router** (DR) and a BDR, the *point-to-point* and the *ptmp* does not.



OSPF Network Types

"This link is working fine, one router reboots and the **DR** *just changes*. When the previous one coming back it *becames* a **BDR.** <u>No problem</u>."



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The Designated Router (DR) and the Backup DR (BDR) are useful when you have **multiple** OSPF routers in the same Layer 2 broadcast domain, to reduce the OSPF traffic.

[The DR maintains a complete topology table of the network and sends updates to the other routers participating in OSPF updates.]



They were designed for scenarios like this one: **more than two OSPF routers** in the same L2 broadcast domain.





Let's now take these details into account and try to look in depths. This is a "typical" OSPF networks between two routers, connected by L2 through L2 devices.





By "default" both of the ETH1 (in Routing / OSPF / Interfaces) will have:

network type=broadcast and priority=1





And as a consequence we will have one **DR** and one **BDR**.



OSPF network



Everything is running fine until something happens





The link goes down. But both of the routers will have the ethernet ports "running", and **THEY BOTH BECOME DR**!!!!



OSPF network



And the question is:

What happen when the L2 Link will be restored?





They will "fight" one against the other, **both are now DR** and no one would like to be BDR anymore.





The result: this **OSPF network is not working**, and usually will require lots of time







The "Engineer" thoughts: "I just rebooted the router and it just worked again"

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The "Engineer" thoughts: "Maybe the OSPF protocol have some bugs" ROUTING & WIRELESS

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The "Engineer" thoughts: "RouterOS does not implement the OSPF protocol well"

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Yes, we should use the proper **network type** for this scenario.

We don't need a DR and a BDR between just **two** routers.







The **point-to-point** is a network type that will run OSPF between **two** routers only and does **not elect** designated router.





The **point-to-point** is a network type that will run OSPF between **two** routers only and does **not elect** designated router.



Using the proper network type between two OSPF routers will save lot of lifes.

Think about that the next time you're doing that!







Is quite easy to use two OSPF paths as: **failover**





ALL the traffic from/to R1 and from/to R2 will go through the PATH 1





In case of failure of the PATH 1 all the traffic from/to R1 and from/to R2 will go through PATH 2.





This is one of the most used OSPF technique in the ISP/WISP industry.





Pros: traffic into the same path

Cons: one "unused" path




Another scenario: load sharing upload/download





Another scenario: load sharing upload/download





In case of failure all the traffic will use the active path only.





Pros: using both paths

Cons: asymmetric traffic





Is it possible to use two different OSPF paths for load sharing and redundancy?



Is it possible to use two different OSPF paths for "constrained" **load sharing** and redundancy?





Is it possible to use two different OSPF paths for "constrained" load sharing **and redundancy**?









The Dijkstra's Algorithm will calculate the sum of the cost of the **output interfaces of each router of the full path**.

Having R1 and R2 with **only one output interface**, like ethernet for example, it is difficult to "steer" the traffic as we would like to.



One solution can be to use multiple physical interfaces.





One solution can be to use multiple physical interfaces.

Like VLANs!





We can use the VLANs to create multiple separate OSPF networks on different "physical output" interfaces.

This way we can steer the traffic as we like to do!

[I've often defined this as: think of "layers"]



The VLANs logic





The results: using both paths





The results: failure of Path 1



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The results: failure of Path 2



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Using the "layers" of VLANs and OSPF networks you can engineer the traffic using ALL the paths ("main" and "backups"), avoid asymmetric traffic and using the same path from/to the same source.





The OSPF "things"



The OSPF "things"

The "OSPF World" is very interesting, this protocol can help you a lot for expanding your network: you have just to "know him well" before enabling 😂

I was not able to explain all the "things" that I planned to: they are too many for our the time slot and location today.



The OSPF "things"

But here at the MUM you can talk with me and the other four MikroTik Trainers of the





Wrap up

✓ I hope you enjoyed my presentation

Iearn or discover at least one thing about OSPF

✓ Use RouterOS as OSPF router!!!!



See you in Riga!







Thank you for listening!

Q & A

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